

Models and Peterson's Algorithm

This is Peterson's algorithm:

```
boolean[] enter = {false, false}; int yield = 0 || 1;

thread t0                                thread t1
1 while (true) {                            10 while (true) {
2 // entry protocol                          11 // entry protocol
3 enter[0] = true;                           12 enter[1] = true;
4 yield = 0;                                  13 yield = 1;
5 await (!enter[1]                             14 await (!enter[0]
    || yield != 0);                             || yield != 1);
6 critical section { ... }                    15 critical section { ... }
7 // exit protocol                             16 // exit protocol
8 enter[0] = false;                           17 enter[1] = false;
9 }                                             18 }
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The successors of $\langle yield = 0, \triangleright 6, enter[0] = T, \triangleright 14, enter[1] = T \rangle$ are:

1. The t_0 successor is $\langle yield = 0, \triangleright 3, enter[0] = F, \triangleright 14, enter[1] = T \rangle$.
2. The t_0 successor is $\langle yield = 0, \triangleright 8, enter[0] = T, \triangleright 14, enter[1] = T \rangle$.
3. The t_1 successor is $\langle yield = 0, \triangleright 6, enter[0] = T, \triangleright 15, enter[1] = T \rangle$.
4. There is no t_1 successor.

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$\langle yield = 0, \triangleright 6, enter[0] = T, \triangleright 15, enter[1] = T \rangle$? Both threads are in their critical sections.

- State $\langle yield = 0, \triangleright 6, enter[0] = T, \triangleright 14, enter[1] = T \rangle$ does not exist.
- There is a bug in Peterson's algorithm.
- The previous slide was wrong.
- State $\langle yield = 0, \triangleright 6, enter[0] = T, \triangleright 14, enter[1] = T \rangle$ is never entered.

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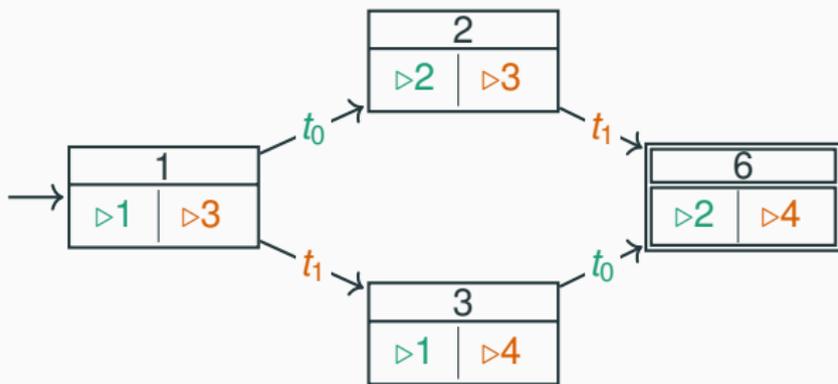
What does Peterson's algorithm achieve?

1. Mutual exclusion using only atomic reads and writes
2. Mutual exclusion and first-come-first-served fairness
3. Mutual exclusion using busy waiting
4. Mutual exclusion using test-and-set operations

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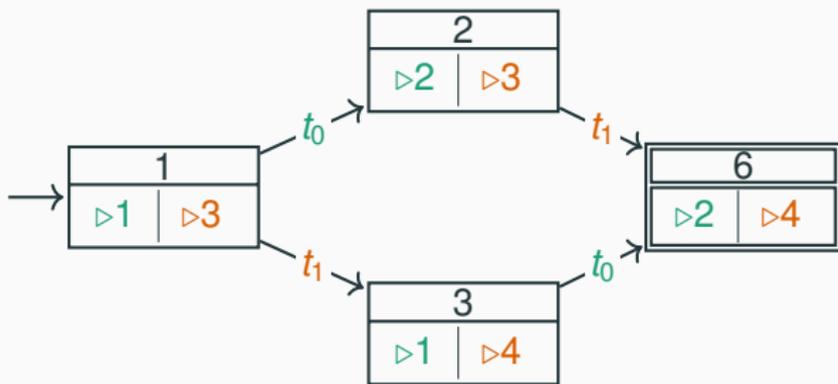
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2. There are no race conditions
3. No starvation can occur, but deadlocks may occur
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